

# CS 3313:

# Foundations of Computing

Spring 2021

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LAs: Marshall Thompson, Oliver Broadrick

<http://gw-cs3313-2021.github.io>

# CS 3313: What is it about ?

- Theoretical foundations of Computer Science
  - It's also a look at history of CS...Concept of computing existed before the first computer was built...
- Two aspects to CS:
  - 1. Fundamental ideas & models underlying computing
  - 2. Engineering techniques HW& SW for design of computing systems
- Theoretical CS had its beginnings in diverse fields
  - Electrical Eng.: studying switching circuits
  - Mathematics: studying foundations of logic
  - Linguists: investigating grammars for natural languages
  - Out of these studies came models that are now central to CS!!

# Questions that drove the Theory of Computing field

- Answer/Ask fundamental (abstract) questions:
  - What is computation –i.e., how do you define what an algorithm is?
  - mathematical models for different types of computing machines ?
    - Why is this an interesting question ?
  - How do you formally define a language
    - Natural language or Programming language

# Course Learning Objectives

- Understand and design different automata (machine) models
  - Apply mathematical reasoning to assert properties of the machines
- Formal models for defining languages – Grammars
- Relationship between languages and automata/machines
- Study fundamental limits of computing (and each machine model), develop ability to determine the computational complexity or solvability of problems
- Further develop the ability to abstract computational problems and mathematical perspective to problem solving
- Above all: *course is about problem solving using Math tools*
  - Mathematical puzzles which abstract “real” computational questions

# Course Schedule - Topics

- Part 1: Regular Languages and Finite State Automata.  
(Weeks 1-5). **familiar territory but now from a math perspective**
  - Finite Automata ....same as Finite State Machines in Hardware!
  - Class of regular languages
  - Properties of regular languages
- Part 2: Context Free Languages and Grammars (weeks 5-10)
  - Formal Grammars – context free grammars, parsing
  - Pushdown Automata – adding simple "memory" to finite state machines
- Part 3: Turing Machines and Computability
  - Turing machine model and Universal Turing machine
  - What is computable ? Proving a problem is not solvable
  - Computational complexity models.

# Instruction team

- Bhagi Narahari
- Grad TAs: Andy Feng and Linsheng Liu
- Undergraduate TAs and LAs:
  - Grant McClearn-UTA (Senior, BS-CS)
  - Oliver Broadrick (Junior, BS-CS)
  - Marshall Thompson (Junior, BS-CS)
- All grading inquiries on homeworks directed to Andy or Linsheng (and then follow up with instructor) via email
  - No posting to piazza
- All discussions on labs and lectures, post to Piazza

# In-class work

- You will learn through in-class activities and exercises most classes (lecture+lab)
  - Must read the material and come to class
- If you are assigned an exercise during class (i.e., an “in class exercise”), you need to complete the exercises by the end of the class/day !
  - We may ask a group/student to present solutions to class

# Course Materials – “confusion will be my epitaph”!

- Course webpage – will have links to syllabus, lecture notes, online resources (and videos when applicable)
  - <http://gw-cs3313-2021.github.io>
- Blackboard will be used for:
  - Synchronous online lectures & labs
  - Homeworks and solutions
  - Reporting grades
- Piazza – for discussions: you’ve used this before...



# Piazza

- you've used this before, so you know the protocols:
  - The purpose of this:
    - to encourage you to ask and answer questions
      - Most of the time, you do better than we do!
      - *Be very careful not to border on plagiarism!*
      - *Don't post your HW solution to the world,*
  - Signup instructions posted on Blackboard
  - Do not expect instant response or substitute Piazza for office hours!
    - Piazza is not manned 24 hours/7 days a week
    - *Sometimes an answer may take more than 24 hours!*
  - Posting on piazza, not the same as telling instructor things
    - E.g. : I'm going to miss the exam!
  - Do **NOT** wait until the last minute to ask for clarifications...
    - The instructors & TAs do NOT plan on spending their weekend checking Piazza!

# Textbooks/Software

- Textbook:
  - Introduction to Formal Languages and Automata, 6<sup>th</sup> edition by Peter Linz (earlier editions will work too), JB Learning.
  - Alternate textbook (a very good book a bit denser): Introduction to the Theory of Computation by Michael Sipser, CEngage publishers.
  - Online notes and resources
- JFLAP – simulator for automata
  - You can install it locally on your laptop
  - Check the tutorial video on the course webpage

# Course Requirements: Grading

- Exams: 50%
  - 3 exams
  - Approximately weeks 5, 10, and Finals week
- Homeworks: 25%
- Quiz and inclass (lab) exercises: 25%
- Grades curved (and scaled as percentage of highest score in class)
  - Check syllabus on website for details.



# Three key concepts

## ■ Languages

- Set of sentences (strings/words) formed from some alphabet
- How do we specify the property?

## ■ Grammars

- A formalism for mathematically defining the properties of a language
- A set of rules for generating the sentences in a formal language

## ■ Automata

- Mathematical model of machines (of different capabilities)
- Formal construct that accepts input, produces output and may have temporary storage and can make decisions

# Formal Languages: Basic Concepts

- Alphabet: set of symbols, i.e.  $\Sigma = \{a, b\}$
- String: finite sequence of symbols from  $\Sigma$ , such as  $v = aba$  and  $w = abaaa$ 
  - Empty string ( $\lambda$ )
  - Substring, prefix, suffix
- Operations on strings:
  - *Concatenation*:  $vw = abaabaaa$
  - *Reverse*:  $w^R = aaaba$
  - *Repetition*:  $v^2 = abaaba$  and  $v^0 = \lambda$
- Length of a string:  $|v| = 3$  and  $|\lambda| = 0$

# Formal Languages: Definitions

- $\Sigma^*$  = set of all strings formed by concatenating zero or more symbols in  $\Sigma$
- $\Sigma^+$  = set of all non-empty strings formed by concatenating symbols in  $\Sigma$

In other words,  $\Sigma^+ = \Sigma^* - \{ \lambda \}$

- A *formal language* is any subset of  $\Sigma^*$

Examples:  $L_1 = \{ a^n b^n : n \geq 0 \}$  and  $L_2 = \{ ab, aa \}$

- A string in a language is also called a *sentence* of the language

# Formal Languages: Set Operations

- A language is a set. Therefore, set operations are defined as usual.
- If  $L_1 = \{ a^n b^n \mid n \geq 0 \}$  and  $L_2 = \{ ab, aa \}$ 
  - Union:  $L_1 \cup L_2 = \{ aa, \lambda, ab, aabb, aaabbb, \dots \}$
  - Intersection:  $L_1 \cap L_2 = \{ ab \}$
  - Difference:  $L_1 - L_2 = \{ \lambda, aabb, aaabbb, \dots \}$
  - Complement:  $L_2 = \Sigma^* - \{ ab, aa \}$
- **Question: Find  $L_2 - L_1$**

# Formal Languages: Other Operations

- New languages can be produced by reversing all strings in a language, concatenating strings from two languages, and concatenating strings from the same language.
- If  $L_1 = \{ a^n b^n \mid n \geq 0 \}$  and  $L_2 = \{ ab, aa \}$ 
  - Reverse:  $L_2^R = \{ ba, aa \}$
  - Concatenation:  $L_1 L_2 = \{ ab, aa, abab, abaa, aabbab, aabbaa, \dots \}$   
The concatenation  $L_2 L_2$  or  $L_2^2 = \{ abab, abaa, aaab, aaaa \}$
  - Star-Closure:  $L_2^* = L_2^0 \cup L_2^1 \cup L_2^2 \cup L_2^3 \cup \dots$
  - Positive Closure:  $L_2^+ = L_2^1 \cup L_2^2 \cup L_2^3 \cup \dots$
- Question: Find  $(L_2 - L_1)^R$



# A better formalism to define language ?

## Some questions

- Set notation works but does not specify a way to generate the words/strings in the language
- Ex: how do you define a syntactically valid C program ?
- Ex: how do you define the construction of a sentence in the English language ?

## More questions....

- Will the piece of C code compile ?
  - How to capture this property

```
/* header info.. */
int foo(int x, int y, char a){
    // body of function
}

int main{
    int i,j,k;
    k= foo(i,j);
    ...
}
```

## More questions....

- What does this (English) sentence mean:  
“Grant made Oliver duck”
  
- What does this (English) sentence mean:  
“Time flies like an arrow.”

# Need for formalism and Math rigor

- We would like to capture precisely, and logically, the properties (problems) in the previous examples
- Ex1: actual and formal parameters (arguments) should match
- Ex 2,3: the language is ambiguous...which is a bad thing in a programming language
  - How do we define, using a mathematical structure, what ambiguity means (in an unambiguous manner 😊 )

# Grammars: Definition

- Precise mechanism to describe the strings in a language

- Definition: A grammar G consists of:

V: a finite set of variable or non-terminal symbols

T: a finite set of terminal symbols (the *alphabet!* )

S: a variable called the start symbol

P: a set of productions (also called production rules)

- Example 1:

$$V = \{ S \}$$

$$T = \{ a, b \}$$

$$P = \{ S \rightarrow aSb, S \rightarrow \lambda \}$$

# Grammars: Derivation of Strings

- Beginning with the start symbol, strings are derived by repeatedly replacing variable symbols with the expression on the right-hand side of any applicable production
- Any applicable production can be used, in arbitrary order, until the string contains no variable symbols.
- Sample derivation using grammar in Example 1:
  - $S \Rightarrow aSb$  (applying first production)
  - $\Rightarrow aaSbb$  (applying first production)
  - $\Rightarrow aabb$  (applying second production)

# The Language Generated by a Grammar

- Definition 2: For a given grammar  $G$ , the language generated by  $G$ ,  $L(G)$ , is the set of all strings derived from the start symbol
- To show a language  $L$  is generated by  $G$ :
  - Show every string in  $L$  can be generated by  $G$
  - Show every string generated by  $L$  is in  $G$

# Grammars for Programming Languages

- The syntax of constructs in a programming language is commonly described with grammars
  - Commonly referred to as Backus-Naur Form (BNF)
- Assume that in a hypothetical programming language,
  - Identifiers consist of digits and the letters a, b, or c
  - Identifiers must begin with a letter
- Productions for a sample grammar:

$\langle \text{id} \rangle \rightarrow \langle \text{letter} \rangle \langle \text{rest} \rangle$

$\langle \text{rest} \rangle \rightarrow \langle \text{letter} \rangle \langle \text{rest} \rangle \mid \langle \text{digit} \rangle \langle \text{rest} \rangle \mid \lambda$

$\langle \text{letter} \rangle \rightarrow a \mid b \mid c$

$\langle \text{digit} \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$



# Grammars for English

- Possible rules for a formal English grammar:
  - <sentence> → <noun phrase> <verb phrase>
  - <verb phrase> → <adverb> <verb> | <verb>
  - <noun phrase> → <pronoun> <noun> | <noun>
- In Ex1: <verb> → duck or <noun> → duck
- In Ex.2: <verb>. → Flies or <noun> → flies

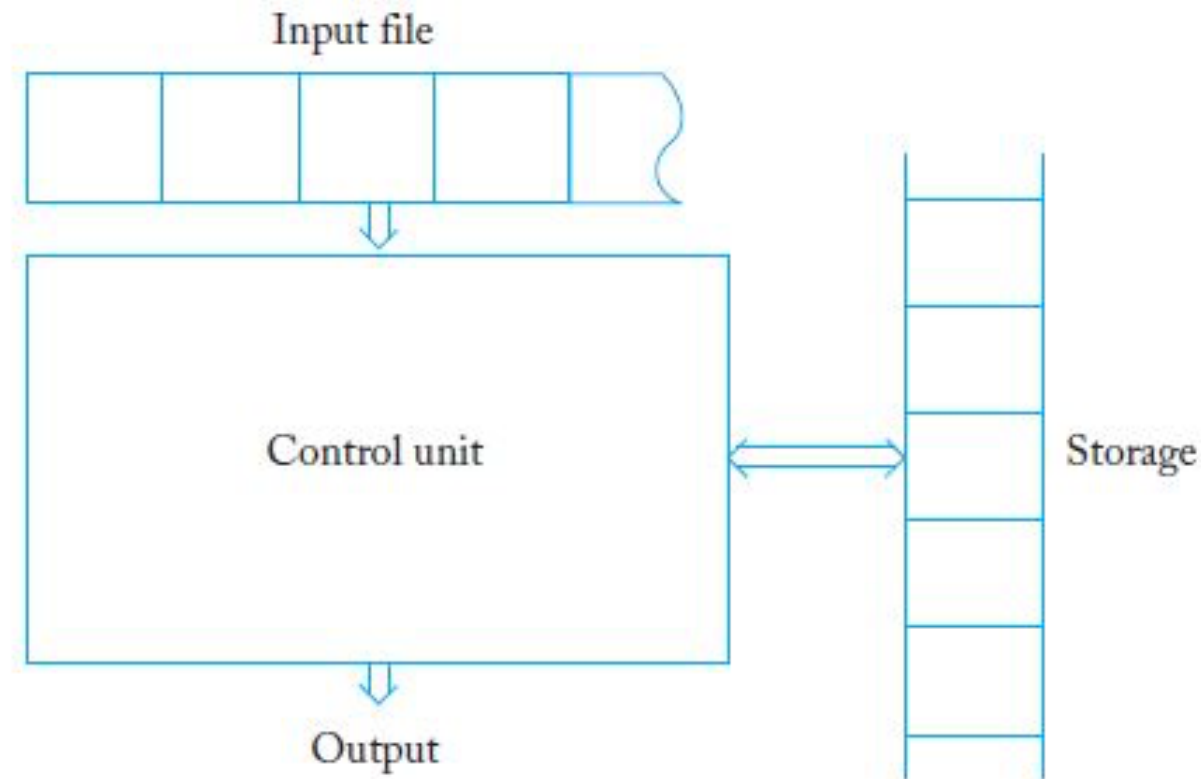
# Concept 3 – modeling computational machines: Automata

- Some “real” machine models:
  - Finite state machines – i.e. sequential circuits
  - von Neumann model of computer architectures
- So why study formal models ?
  - To investigate the computational power of these machines  
what problems can they solve ?
- Question: Can we build a compiler using just a finite state machine ?
  - advantages: simpler design
- Bigger question: what problems are solvable by a von Neumann computer ?

# Automata

- An automaton is an abstract model of a digital computing device
- An automaton consists of
  - An input mechanism
  - A control unit
  - Possibly, a storage mechanism
  - Possibly, an output mechanism
- Control unit can be in any number of internal *states*, as determined by a *next-state* or *transition function*.
- There are a **finite number of states**
  - *Why ?*

# Illustration of a General Automaton



Note: this is one model of the “type” of external storage  
The storage model changes based on the automata model  
Question: Is there storage in a finite state machine ?

# Automata we will study

- Finite Automata (Deterministic & Non-deterministic)
  - These model Finite State Machines
- Pushdown automata
  - Add the simplest form of memory to a Finite state machine
- Turing Machines
  - These model today's computers in terms of computational ability
- Questions to ask: What is the "power" (limit) of each of these machine models ?
  - What types of problems can be solved by a finite automaton, Turing machine etc.
- Link b/w Languages and Automata: view automata as "acceptors" of input strings
  - What languages are accepted by the automata model

# Limits of Computation (of automata): Questions

- Can we use a finite state machine to build a compiler ?
- Can we use a pushdown automaton to parse a programming language ?
- Can we design a compiler that will determine for any program, whether the program halts on all inputs
  - Or, will the compiler detect any bugs in the program ?
  - Or, are two programs equivalent ? (do they compute the same function)
- To answer these questions, via proofs, we need mathematical models for these types of machines!

# Mathematical Rigor

- This is a theoretical course and requires formalisms and formal (math) methods.....
  - Unless we ask for a JFLAP simulation, your answers need to be grounded in Maths (i.e., no multi-page discussions with no mathematical logic!)

## Next...

- We start with Finite Automata
- These are the 'same' as Finite State machines covered in Computer Architecture
  - But we model them as acceptors of languages:  
Input is a string and the Automata determines if the string is in the language

Chapter 2 of Textbook